



Lambdas

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First-Class Functions

- In Racket, functions are **first-class** values
- Can be bound to vars, returned from fns, etc..
- Languages w/ functions as values are functional

Lambdas (in Racket)

- (lambda (x0 x1 ...) body)
 - Anonymous function: bind x0, ... in body
 - Can appear at any callsite (just like an identifier)

```
(define f (lambda (x) x))
(define (double g)
  (lambda (x) (g (g x))))
```

Exercise

```
(define f (lambda (x) x))
(define (double g)
  (lambda (x) (g (g x))))
```

Evaluate the following expressions:

- (f 1)
- ((double f) 42)
- ((double (lambda (x) (* x 2))) 2)



Write a function, (foo f), that:

- Accepts a function f, maps ints to ints
- ((foo f) x) = (f |x|), |x| is abs. value of x

Textual Reduction of Lambdas

- Previously, we assumed **environment** of definitions
- Instead, can think of lambdas as primitive
- Environment maps identifiers to lambdas

```
(define (f x) x)
;; equiv
(define f (lambda (x) x))
```

Textual Reduction of Lambdas

• After reducing all args to values, substitute (into the body) the actual arguments in place of the formal arguments.

```
((lambda (x y) x) (+ 1 1) 3)
=> ((lambda (x y) x) 2 3)
=> 2
```



Use textual reduction to reduce the following:

Hint: remember, in applicative order we always evaluate the leftmost, innermost application. In other words, we process (e0 e1 ...) by reducing e0 ... to values in order, then applying.

Exercise

Use textual reduction to reduce the following:

If this sounds complicated, you would be right to just think about it as "left to right"

Languages w/o First-Class Functions

- In modern times, somewhat hard to imagine
- C is a good example: procedural but **not** functional
- C callsites: quasi-functional behavior via fn pointers
 - But not really: C doesn't have closures