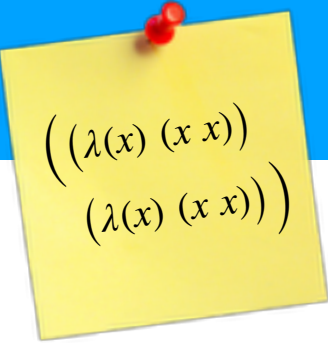


# First-class and Higher-order Functions

CIS 352 — Spring 2020

# Example

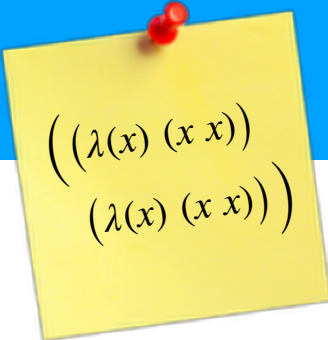


$((\lambda(x) (x x))$   
 $(\lambda(x) (x x)))$

Squaring every element of a list

```
(define (square-list-values lst)
  (if (null? lst)
      '()
      (cons (* (car lst) (car lst))
            (square-list-values (cdr lst)))))
```

# Example



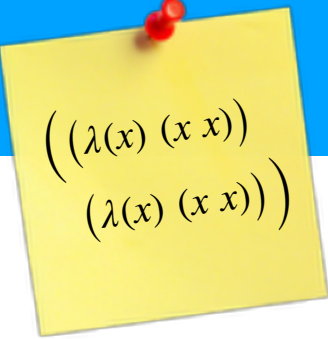
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(define (square-list-values lst)
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      (cons (* (car lst) (car lst))
            (square-list-values (cdr lst)))))
```

Defines **base case**

# Example



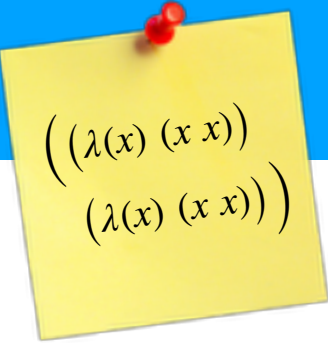
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(define (square-list-values lst)
  (if (null? lst)
      '()
      (cons (* (car lst) (car lst))
            (square-list-values (cdr lst)))))
```

**Recursive case** first computes the square of (car lst)

# Example



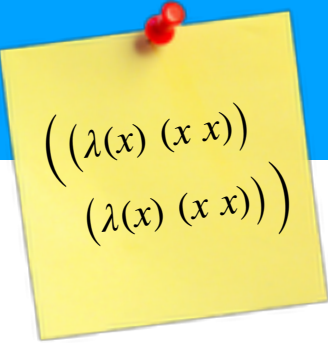
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Squaring every element of a list

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(define (square-list-values lst)
  (if (null? lst)
      '()
      (cons (* (car lst) (car lst))
            (square-list-values (cdr lst)))))
```

**Recursive case** next recurs on the list's tail (cdr lst)

# Example



$((\lambda(x) (x x))$   
 $(\lambda(x) (x x)))$

Squaring every element of a list

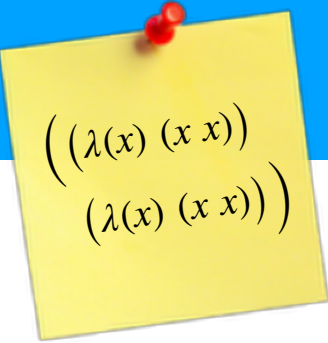
```
(define (square-list-values lst)
  (if (null? lst)
      '()
      (cons (* (car lst) (car lst))
            (square-list-values (cdr lst)))))
```

**Recursive case** finally extends the new tail list

# Anonymous Functions

- Like many languages (e.g., JS, Python, Ruby), Racket allows anonymous functions to be defined and treated as **values**.
  - `(lambda (args ...) body)` ; returns a function as a value
  - E.g., `(lambda (x) (* x x))` ; returns a square function
  - When a language permits functions to be treated as any other value may be treated (passed to other functions, bound to variables, stored in a list, etc), such functions are called **first-class** functions.
- Actually, all functions are anonymous—these are not special.
  - `(define (id x) x) == (define id (lambda (x) x))`

# Example



$((\lambda(x) (x x)))$   
 $(\lambda(x) (x x))$

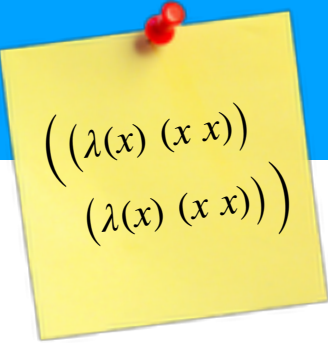
Squaring every element of a list

```
(define (map f lst)
  (if (null? lst)
      '()
      (cons (f (car lst))
            (map f (cdr lst)))))
```

```
(define (square-list-values lst)
  (map (lambda (x) (* x x)) lst))
```




# Example



$((\lambda(x) (x x))$   
 $(\lambda(x) (x x)))$

Squaring every element of a list

```
(define (map f lst)
  (if (null? lst)
      '()
      (cons (f (car lst))
            (map f (cdr lst)))))
```



**map** takes a  
(unary) function  
and list

```
(define (square-list-values lst)
  (map (lambda (x) (* x x)) lst))
```

# Example

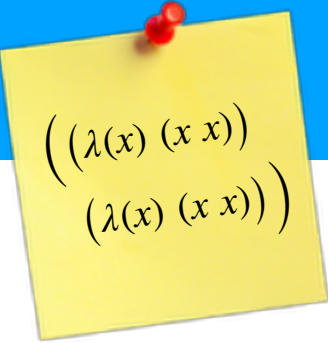
$((\lambda(x) (x x))$   
 $(\lambda(x) (x x)))$

Squaring every element

```
(define (map f lst)
  (if (null? lst)
      '()
      (cons (f (car lst))
            (map f (cdr lst)))))
```

Works essentially as  
**square-list-values**  
except each new  
value is **(f (car lst))**


```
(define (square-list-values lst)
  (map (lambda (x) (* x x)) lst))
```



$(\lambda(x) (x x))$   
 $(\lambda(x) (x x))$

Squaring every element of a list

Now we may define  
**square-list-values**  
(in one line) in terms of our  
(highly-reusable) **map** component

  
`(define (square-list-values lst)  
 (map (lambda (x) (* x x)) lst))`

# Exercise



What is the return value of the following expression?

```
(let ([f (lambda (a) (* a a a))])  
  (let ([g add1])  
    (let ([h f])  
      (g (h 5))))))
```

# Exercise



What is the return value of the following expression?

```
(let ([f (lambda (a) (* a a a))])  
  (let ([g add1])  
    (let ([h f])  
      (g (h 5)))))
```

**Answer: 126**

# Exercise



What is the return value of the following expression?

```
(let ([tw (lambda (f x) (f (f x)))]  
      [th (lambda (f x) (f (f (f x)))]])  
  (let ([f add1])  
    (tw (lambda (x) (th add1 x)) 0))))
```

# Exercise



What is the return value of the following expression?

```
(let ([tw (lambda (f x) (f (f x)))]  
      [th (lambda (f x) (f (f (f x)))]])  
  (let ([f add1])  
    (tw (lambda (x) (th add1 x)) 0))))
```

**Answer: 6**

# Higher-order functions

- Languages with first-class functions also have higher-order (HO) functions and are called **higher-order languages** (HOL).
  - A **higher-order function** is a function over functions: a function that takes a function as input, returns a function as output, or both.
- Common higher-order functions include `map`, `foldl`, `foldr`, `filter`, `andmap`, `ormap`, etc...
  - `foldl/foldr` walks a list and uses a function to reduce it
  - `map` walks a list to turn every `x` into `(f x)` for parameter `f`
  - `andmap/ormap` lift a predicate (`param`) to a list predicat





Write an implementation of `andmap`, such that:

```
> (andmap list? '((1 2) () (3)))
```

```
#t
```

```
> (andmap list? '((1 . 2) ()))
```

```
#f
```

```
> (andmap list? '(1 2 3))
```

```
#f
```



Double-check: does your implementation **short-circuit**? What does your implementation give for:

```
> (andmap list? '())
```



## Answer:

```
(define andmap
  (lambda (p? lst)
    (if (null? lst)
        #t
        (and (p? (car lst))
              (andmap p? (cdr lst))))))
```

A predicate  $p?$   
*trivially* holds for all  
elements of '()



## Answer:

```
(define andmap
  (lambda (p? lst)
    (if (null? lst)
        #t
        (and (p? (car lst))
              (andmap p? (cdr lst))))))
```



This short-circuits because (and ...) does!



Another definition, without using (and ...):

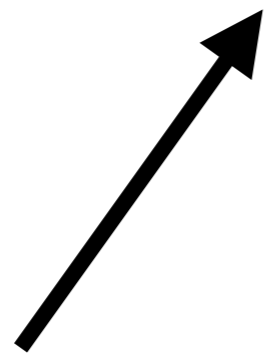
```
(define (andmap p? lst)
  (if (null? lst)
      #t
      (if (p? (car lst))
          (andmap p? (cdr lst))
          #f)))
```

Use an if to check the next element. If the test fails, short-circuit and return #f, otherwise recur.



Yet another definition, using a fold:

```
(define (andmap p? lst)
  (foldl (λ (elem b)
          (and b (p? elem)))
        #t
        lst))
```



fold over the list, accumulating a single boolean:  
at each step, conjoin this bool with `(p? elem)`



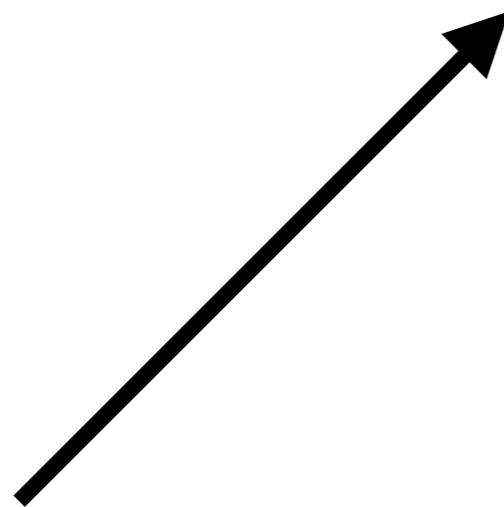
Write an implementation of map, using a fold:

```
> (map add1 '(1 2 3))  
'(2 3 4)
```



## Answer:

```
(define (map f lst)
  (foldr (lambda (x tail)
          (cons (f x) tail))
        '()
        lst))
```



Fold over the list from right-to-left, accumulating an updated tail of the list, replacing each  $x$  with  $(f x)$



# Free variables

- A variable  $x$  is called **free** in expression  $e$ , if there exists a reference to  $x$  within  $e$  whose definition is not also within  $e$ .
  - E.g.,  $x$  is free in **(let ([y 3]) (+ x y))**, but  $y$  is not.
  - E.g.,  $x$  is free in **(list x y z)**; so are  $y$ ,  $z$ , and **list**!
  - Expressions with no free variables are valid programs!
- A function with no free variables is called a **combinator**.
  - E.g., **(lambda (x) (\* x x))** or **(λ (f x) (f (f x)))**
  - Combinators are stand-alone, reusable components
- Functions with free variables, save their values! (*More soon*)

# Exercise



What are the free variables of the high-lit expression?

```
(let ([f (lambda (x)
          (lambda (y)
            (+ x y)))]])
  (let ([g (f 2)])
    (g 3)))
```

# Exercise



What are the free variables of the high-lit expression?

```
(let ([f (lambda (x)
          (lambda (y)
            (+ x y)))]])
  (let ([g (f 2)])
    (g 3)))
```

**Answer: { f }**

# Exercise



What are the free variables of the high-lit expression?

```
(let ([h (λ (x) (+ 3 x))])  
  (let ([g (λ (x y) (* x y y))])  
    (h (g 3 4))))
```

# Exercise



What are the free variables of the high-lit expression?

```
(let ([h (λ (x) (+ 3 x))])  
  (let ([g (λ (x y) (* x y y))])  
    (h (g 3 4))))
```

**Answer: { h, \* }**

# Exercise



What are the free variables of the high-lit expression?

```
(lambda (x)
  (lambda (y)
    ((lambda (x z) (- x z)) x y)))
```

# Exercise



What are the free variables of the high-lit expression?

```
(lambda (x)
  (lambda (y)
    ((lambda (x z) (- x z)) x y)))
```

**Answer: { x, y, - }**

# Currying

- Using higher-order functions, it is always possible to encode a k-ary function as a set of unary functions via **currying**:
  - Invented by Frege; popularized by Schönfinkel, Curry
  - A function **(define twice (λ (f x) (f (f x))))** is curried as two nested functions:  
**(define twice (λ (f) (λ (x) (f (f x)))))**  
and to apply the function we call it twice  
**((twice add1) 0)**
  - The first call binds **f** to **add1** and returns a function that *saves / remembers* this value for **f**.
  - The second call binds **x** and returns **(f (f x))**





Define a curried version of the slope function:

```
(define (slope x0 y0 x1 y1)
  (/ (- y1 y0) (- x1 x0)))
```

```
> (slope 1 1 5 9)
2
```



## Answer:

```
(define (slope x0)
  (lambda (y0)
    (lambda (x1)
      (lambda (y1)
        (/ (- y1 y0) (- x1 x0)))))))
```

```
> (((slope 1) 1) 5) 9)
2
```